

# A PATH COUNTING PROBLEM IN DIGRAPHS

KAREL ZIKAN and EDWARD SCHMEICHEL  
San Jose State University, San Jose, CA 95192  
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## 1. INTRODUCTION

In this paper, we consider only directed graphs without loops or multiple edges. Our terminology and notation will be standard except as noted. A good reference for any undefined terms is [1].

Our main goal is to determine the maximum possible number of directed paths between a pair of vertices in an acyclic digraph with  $m$  edges (and any number of vertices). Denoting this maximum possible number by  $N(m)$ , we will establish that

$$N(m) = \begin{cases} F_{(m+1)/2} & \text{for } m \text{ odd} \\ 1 & \text{for } m = 2 \\ 2F_{(m/2)-1} & \text{for } m \geq 4 \text{ and even} \end{cases}$$

where  $F$  satisfies the recurrence relation

$$F_k = F_{k-1} + F_{k-2}, \quad F_1 = 1, \quad F_2 = 2.$$

The actual proof of this formula will be preceded by a sequence of five easy lemmas.

We then conclude with a brief discussion of the following related question: Given a positive integer  $k$ , what is the least number of edges in an acyclic digraph having *exactly*  $k$  directed paths between a pair of vertices.

## 2. PROOFS OF THE LEMMAS AND MAIN RESULT

### Lemma 1

Let  $D$  be an acyclic digraph. Then  $D$  contains vertices  $a$  and  $z$  such that indegree  $a$  = outdegree  $z$  = 0. (We call  $a$  and  $z$ , respectively, a source and a sink of  $D$ .)

Proof: Let  $x \in V(D)$ . Consider a longest path directed away from  $x$ , say from  $x$  to  $z$ . Then outdegree  $z$  = 0 (since any edge leaving  $z$  would yield either a longer directed path away from  $x$  or a directed cycle in  $D$ ).

The proof that indegree  $a$  = 0 for some  $a \in V(D)$  is entirely analogous. ■

### Lemma 2

Let  $D$  be an acyclic digraph. Then the vertices of  $D$  can be ordered, say  $1, 2, \dots, n$ , such that every edge in  $D$  is of the form  $(i, j)$ , where  $i < j$ .

Proof: We proceed by induction on  $n = |V(D)|$ . The result is trivially true for  $n = 2$ . For the induction step, choose any  $z \in V(D)$  with outdegree  $z$  = 0 (one exists by Lemma 1), and consider the digraph  $D - z$ . By the induction hypothesis, the vertices of  $D - z$  can be ordered, say  $1, 2, \dots, n - 1$ ,

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in the manner described. If we let  $z$  be the  $n^{\text{th}}$  vertex, we have the desired ordering of  $V(D)$ . ■

In what follows, we assume  $D$  is an acyclic digraph with vertices ordered  $1, 2, \dots, n$  such that every edge of  $D$  is of the form  $(i, j)$ , where  $i < j$ .

For any  $x \in V(D)$ , let  $p_D(x)$  denote the number of directed paths from 1 to  $x$  in  $D$ . [When  $D$  is clear from context, we will use just  $p(x)$  for this number.]

### Lemma 3

Suppose  $D$  has a set of vertices  $S = \{i < \dots < j < k\}$ , with  $1 < i < k \leq n$ , which induces a tournament (i.e., a digraph with every pair of vertices joined by precisely one edge). Then

$$p(k) \geq p(i) + \dots + p(j).$$

Proof: For each  $x \in S$ , let  $P(x)$  denote the set of directed paths from 1 to  $x$ . If  $x \neq k$ , let  $P'(x)$  denote the set of directed paths from 1 to  $k$  obtained by taking a path from 1 to  $x$  together with the edges  $(x, k)$ . Then, clearly,

$$P'(i) \cup \dots \cup P'(j) \subseteq P(k),$$

and the sets on the left side are disjoint. Since

$$|P'(x)| = |P(x)| = p(x),$$

it follows at once that

$$p(i) + \dots + p(j) \leq p(k). \quad \blacksquare$$

Let  $N(m)$  denote the maximum possible number of directed paths between two vertices of an acyclic digraph with  $m$  edges. Certainly  $N(m)$  is a nondecreasing function of  $m$ . Let us call an acyclic digraph on  $m$  edges having precisely  $N(m)$  directed paths between some pair of vertices a *path maximum  $m$ -graph*. It is easily seen that there will be a path maximum  $m$ -graph  $D$  with the vertices ordered as in Lemma 2 such that 1 and  $n$  are joined by precisely  $N(m)$  directed paths, and 1 (resp.,  $n$ ) is the unique source (resp., sink) in  $D$ . We will assume this property for the path maximum  $m$ -graphs we consider in what follows.

### Lemma 4

There exists a path maximum  $m$ -graph  $D$  in which

$$\{x \in V(D) \mid (x, n) \in E(D)\}$$

(i.e., the predecessors of  $n$  in  $D$ ) induce a tournament.

Proof: Otherwise, let  $i, j$  be two predecessors of  $n$  (with say  $i < j$ ) such that  $(i, j) \notin E(D)$ . Form the digraph

$$D' = D - (i, n) + (i, j).$$

To each directed path in  $D$  from 1 to  $n$  containing the edge  $(i, n)$  there corresponds uniquely a directed path in  $D'$  from 1 to  $n$  containing the edges  $(i, j)$  and  $(j, n)$ . Hence,  $p_{D'}(n) \geq p_D(n)$ , and so  $D'$  is also a path maximum  $m$ -graph in which  $n$  has one less predecessor than in  $D$ . We simply iterate this procedure until we obtain a path maximum  $m$ -graph with the desired properties. ■

### Lemma 5

If  $m \geq 3$ , there exists a path maximum  $m$ -graph in which  $n$  has indegree 2.

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Proof: Let  $D$  be a path maximum  $m$ -graph in which the predecessors of  $n$  (ordered say  $1 < \dots < j < k$ ) induce a tournament. By Lemma 3,

$$p(k) \geq p(i) + \dots + p(j).$$

Hence,

$$2p(k) \geq p(i) + \dots + p(j) + p(k) = p(n) = N(m).$$

If indegree  $n \geq 3$ , we can construct a new acyclic digraph  $D'$  with  $m$  edges, as shown in Figure 1. Note that

$$p_{D'}(n') = 2p(k) \geq N(m),$$

and hence  $D'$  is also a path maximum  $m$ -graph. But indegree  $_{D'} n' = 2$ , and the proof is complete. ■

(indegree  $n$ ) - 1 edges

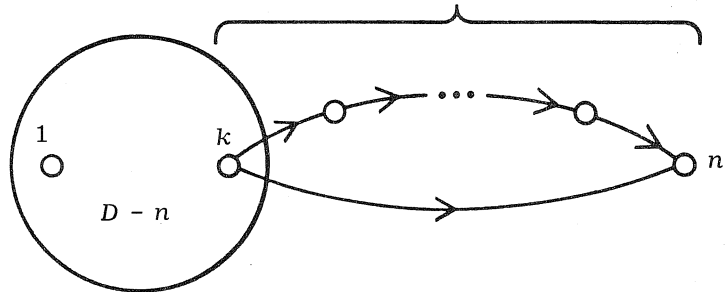


Figure 1. The Digraph  $D'$

We now state and prove our main result.

Theorem

Let  $m$  be a positive integer. Then

$$N(m) = \begin{cases} F_{(m+1)/2} & \text{for } m \text{ odd} \\ 1 & \text{for } m = 2 \\ 2F_{(m/2)-1} & \text{for } m \geq 4 \text{ and even} \end{cases}$$

where  $F_k$  is the Fibonacci number satisfying  $F_k = F_{k-1} + F_{k-2}$ ,  $F_1 = 1$ ,  $F_2 = 2$ .

Proof: It is readily verified that

$$N(1) = N(2) = 1, N(3) = N(4) = 2, N(5) = 3, N(6) = 4,$$

and so the formula holds for  $m \geq 6$ . We thus proceed by induction on  $m \geq 7$ .

Since the digraphs in Figure 2 contain  $m$  edges, and have as many dipaths from 1 to  $n$  as the number specified in the formula, it suffices to show the numbers in the formula are upper bounds for  $N(m)$ .

By Lemma 5 there is a path maximum  $m$ -graph  $D$  in which the indegree of  $n$  is 2. Let  $x, y$  denote the predecessors of  $n$  in  $D$ , with say  $x < y$ . We then have precisely three possibilities:

- (i)  $(x, y) \notin E(D)$  (Using the construction in the proof of Lemma 4, we could obtain a path maximum  $m$ -graph in which  $n$  has indegree 1.)

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- (ii)  $(x, y) \in E(D)$ , and  $x$  is the only predecessor of  $y$ .
- (iii)  $(x, y) \in E(D)$ , and  $x$  is not the only predecessor of  $y$ .

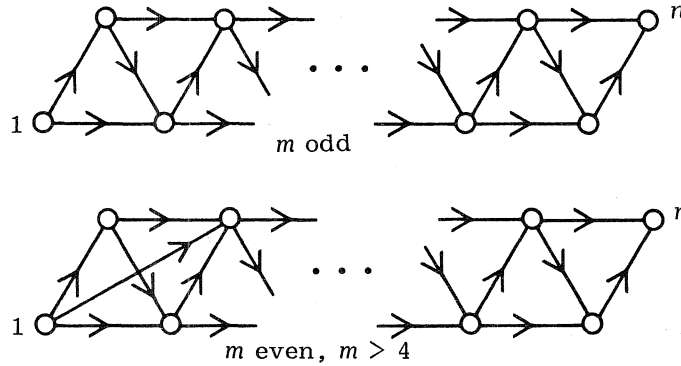


Figure 2. Path Maximum  $m$ -Graphs

By considering the maximum possible number of dipaths from the source to  $x$  and  $y$  in cases (i), (ii), and (iii), respectively, we get

$$N(m) \leq \max\{N(m-1), 2N(m-3), N(m-2) + N(m-4)\}.$$

Using the induction hypothesis, and the fact that  $m \geq 7$ , we obtain

$$N(m) \leq \begin{cases} \max\{2F_{(m-3)/2}, 4F_{(m-5)/2}, F_{(m-1)/2} + F_{(m-3)/2}\} = F_{(m+1)/2}, & \text{if } m \text{ odd,} \\ \max\{F_{(m/2)}, 2F_{(m/2)-1}, 2F_{(m/2)-2} + 2F_{(m/2)-3}\} = 2F_{(m/2)-1}, & \text{if } m \text{ even.} \end{cases}$$

The inductive step, and hence the proof of the theorem, are now complete. ■

### 3. A RELATED PROBLEM

The authors have also considered the following problem: Given a positive integer  $k$ , what is the least number of edges in an acyclic digraph having *exactly*  $k$  paths between some pair of vertices? Noting the  $N(m)$  is nondecreasing in  $m$ , it seems reasonable to conjecture that if  $N(m-1) < k \leq N(m)$ , then  $m$  is the minimum number of edges required. This conjecture is indeed true for  $k \leq 32$ . However,  $N(14) < 33 < N(15)$ , and we have shown that at least 16 edges are needed in any digraph having exactly 33 directed paths between a pair of vertices. Although it appears that a complete solution to this problem may be very difficult, we have the following conjecture to offer:

Conjecture: Let  $k_n$  be the smallest integer such that  $N(m-1) < k_n < N(m)$ , but at least  $m+n$  edges are needed in any digraph with precisely  $k_n$  directed paths between a pair of vertices. Then  $k_n$  satisfies the recurrence relation  $k_n = 34k_{n-1} + 21$ ,  $k_1 = 33$ .

#### REFERENCE

1. M. Behzad, G. Chartrand, & L. Lesniak-Foster. *Graphs and Digraphs*. Boston, Mass.: Prindle, Weber and Schmidt, 1979.

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